### 6.4 FORD-FULKERSON DEMO


click to begin demo
, Ford-Fulkerson algorithm

## Ford-Fulkerson algorithm

Initialization. Start with 0 flow.
initialization

Idea: increase flow along augmenting paths

Augmenting path. Find an undirected path from $s$ to $t$ such that:

- Can increase flow on forward edges (not full).
- Can decrease flow on backward edge (not empty).
$1^{\text {st }}$ augmenting path


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Termination. All paths from $s$ to $t$ are blocked by either a

- Full forward edge.
- Empty backward edge.
no more augmenting paths

, computing a min cut

Computing a mincut from a maxflow

To compute mincut $(A, B)$ from maxflow $f$ :

- By augmenting path theorem, no augmenting paths with respect to $f$.
- Compute $A=$ set of vertices connected to $s$ by an undirected path with no full forward or empty backward edges.


Computing a mincut from a maxflow

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